ON TELESCOPING LINEAR EVALUATION FUNCTIONS

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1. THE RESULT

Consider some finite set X, for instance the set of all legal chess positions. A real-valued function $f: X \to \mathbb{R}$ is called an *evaluation function* on X. In chess, one conventionally prefers evaluation functions that assume small values for positions unfavourable for White and large values for positions favourable for White.

Given *m* arbitrary evaluation functions $f_1, ..., f_m$ on *X*, and real numbers $c_1, ..., c_m$, one can construct a new evaluation function *f* on *X* by taking

$$f(x) = \sum_{i=1}^{m} c_i f_i(x)$$
 for all positions x.

f is called a *linear combination* of $f_1, ..., f_m$. Sometimes the f_i are called elementary evaluation functions. A linear combination f is called *convex* if the coefficients c_i satisfy the following conditions:

$$c_i \ge 0$$
 for all i and $\sum_{i=1}^m c_i = 1$.

In Althöfer (1993), I have proved the following theorem.

Theorem: Let X be a finite set with n elements. Let f_i , ..., f_m be evaluation functions on X with $0 \le f_i(x) \le 1$ for i = 1, ..., m and all $x \in X$. Let $f = \sum_{i=1}^m c_i f_i$ be some convex combination of the f_i . Choose any constant $\varepsilon > 0$ and let $k = \left\lfloor \frac{\ln 2n}{2\varepsilon^2} \right\rfloor$. Then there exists another convex combination $g = \sum_{i=1}^m d_i f_i$ of the same f_i with the following properties: (i) $|g(x) - f(x)| \le \varepsilon$ for all $x \in X$

and

(ii) at most k coefficients d_i differ from zero. Moreover, the $d_i > 0$ are of the form $d_i = k_i/k$ with natural numbers k_i .

[z] denotes the ceiling, i.e., the smallest integer not smaller than the real number z.

Interpretation: (i) states that g is an ε -approximation of f; (ii) states that the number of terms only grows logarithmically with the number of elements in the set X and that the d_i are rational numbers.

The operation of finding k coefficients to represent a function which possibly required many more coefficients is known as telescoping.

The theorem states the existence and does not describe a construction. However, a known result not proved here is that for every fixed $\varepsilon > 0$ g can be constructed by a probabilistic algorithm in expected time O (n ln n) (Althöfer, 1993). This means that the problem, being $O(n \ln n)$, is as hard as sorting which is also $O(n \ln n)$. ln n) (Knuth, 1973). Moreover, halving ε increases computation time by a factor of 4. Also note that the construction assumes that the simple f_i (*i*=1, ..., *m*), the elementary evaluation functions, are given, which regrettably is not the case.

2. APPLICATIONS TO CHESS AND OTHER GAMES

It has been estimated that there are approximately 10⁵³ many legal chess positions (Allis, Van den Herik and Herschberg, 1991). For convenience, in all the following examples the elementary evaluation functions f_i are assumed to satisfy $0 \le f_i(x) \le 1$ for all positions $x \in X$.

Example 1

Consider the set X of all legal chess positions and a convex combination $f = \sum_{i=1}^{m} c_i f_i$ of elementary evaluation functions f_i which satisfies the following strong separation properties:

(i)	$0 \le f(x) < \frac{1}{7}$	if x is a loss for White
(ii)	$\frac{3}{7} < f(x) < \frac{4}{7}$	if x is a drawn position
(iii)	$6/_{7} < f(x) \le 1$	if x is a win for White

Then there exists another (simple) convex combination $g = \sum_{i=1}^{m} d_i f_i$ of the same f_i which satisfies the weaker

separation properties

(i)	$0 \le g(x) < \frac{2}{7}$	if x is a loss for White
(ii)	$\frac{2}{7} < g(x) < \frac{5}{7}$	if x is a drawn position
(iii)	${}^{5}/_{7} < g(x) \leq 1$	if x is a win for White
and		

(iv) at most 3007 coefficients d_i differ from zero.

Proof: Put cardinality $(X) = 10^{53}$, $\varepsilon = \frac{1}{7}$, and apply the theorem.

Example 2

Consider some chess endgame (for instance with 5 pieces) with $10^8 = 100$ million legal positions x. Assume that there is some convex combination $f = \sum_{i=1}^{m} c_i f_i$ which satisfies the strong separation properties

(i) $0 \leq f(x) < \frac{1}{4}$ if x is not a win for White $\frac{3}{4} < f(x) \le 1$ (ii) if x is a win for White

Then there exists another (simple) convex combination $g = \sum_{i=1}^{m} d_i f_i$ of the same f_i which satisfies the weaker separation properties

- $\begin{array}{l} 0 \leq g(x) < \frac{1}{2} \\ \frac{1}{2} < g(x) \leq 1 \end{array}$ (i) if x is not a win for White
- if x is a win for White (ii)

and

(iii) at most 153 coefficients d_i differ from zero.

Proof: Put card $(X) = 10^8$, $\varepsilon = \frac{1}{4}$, and apply the theorem.

Example 3

Consider the chess endgame KBNNKR (Stiller and the Editors, 1991). There are fewer than $n = 32 \times 63 \times 62 \times 61 \times 60 \times 59 \times 2$ legal positions in this endgame. Assume than there is some convex combination f =

 $\sum_{i=1}^{m} c_i f_i$ which satisfies the strong separation properties

(i) $0 \le f(x) < \frac{1}{4}$ if x is not a win for White (ii) $\frac{3}{4} < f(x) \le 1$ if x is a win for White.

Then there exists another (simple) convex combination $g = \sum_{i=1}^{m} d_i f_i$ of the same f_i which satisfies the weaker separation properties

separation properties

(i) $0 \le g(x) < \frac{1}{2}$ if x is not a win for White (ii) $\frac{1}{2} < g(x) \le 1$ if x is a win for White and

(iii) at most 204 coefficients d_i differ from zero.

Proof: Put card $(X) = 32 \times 63 \times 62 \times 61 \times 60 \times 59 \times 2 \approx 5.4 \ 10^{10}$, $\varepsilon = \frac{1}{4}$ and apply the theorem.

Example 4

Consider the game of Go on a 19 × 19 board. Exploiting symmetries, there are fewer than 3^{361} legal positions. Assume that there is some convex combination $f = \sum_{i=1}^{m} c_i f_i$ of elementary evaluation functions f_i which satisfies the strong separation properties

(i) $0 \le f(x) < \frac{1}{4}$ if x is not a win for Black (ii) $\frac{3}{4} < f(x) \le 1$ if x is a win for Black

Then there exists another (simple) convex combination $g = \sum_{i=1}^{m} d_i f_i$ of the same f_i which satisfies the weaker

separation properties

(i)	$0 \le g(x) < \frac{1}{2}$	if x is not a win for Black
(ii)	$\frac{1}{2} < g(x) \leq 1$	if x is a win for Black
and		

(iii) at most 3179 coefficients d_i differ from zero.

Proof: Put card $(X) = 3^{361}$, $\varepsilon = 1/4$, and apply the theorem.

Of course the theorem can be applied also to all other finite games, for instance those investigated in Allis et al. (1991).

3. CONCLUSIONS

- (i) Theorem 1 states: "if there is a convex combination $f = \sum c_i f_i$... then ...". It does not answer the question whether such an f is so expressible; only if it is, the construction of $g = \sum d_i f_i$ from $f = \sum c_i f_i$ becomes possible. In other words, without knowing the c_i we may not construct the d_i .
- (ii) Theorem 1 may be interpreted in the context of the *search-versus-knowledge* discussion: convex combinations of elementary evaluation functions with very many terms may often have shorter counterparts in which the deterioration of quality may be controlled.

(iii) Of course one may argue that convex combinations are often not the best way to use (large) sets of elementary evaluation functions.

4. **REFERENCES**

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REVIEW

SECRETS OF ROOK ENDINGS

by John Nunn

London, UK

B.T. Batsford Ltd., 1992 London, UK 320 pages ISBN 0-7134-7164-6

Reviewed by H.J. van den Herik and I.S. Herschberg

The very title of John Nunn's most recent book reads more like a novel's and indeed it is a novel kind of book. The novelty is in its utter perfection and unassailable solidity. It is deeply satisfying to find the exclamation mark used with a precise, stated definition, instead of depending on a commentator's passing fancy. The query, too, now is perfectly objective. These are just two minor delights in a delightful major work showing "the almost inexhaustible nature of chess".

It must not be thought that this book is merely a digest of Ken Thompson's KRPKR database (Thompson, 1990), which later was made public on a first volume of a series of CD-ROMs (see Thompson, 1991). Lars Rasmussen was an important intermediary, providing a user-friendly interface which rendered the ultracompact database grandmaster-friendly. Needless to say, that the systematics of this book leaves nothing to be desired: its main subdivision is into Chapters, one each for the pawn files a, b, c, and d. Within each chapter, the subdivision is by pawn rank. Within that ranking, human terminology induces a further grouping according to whether the King or Rook is in front of or behind the Pawn, defends it from the side, and a miscellaneous section.

The book is enlivened by no fewer than 534 diagrams on 320 pages. Rolf Schlösser's Postscript chess fonts are uniform in text and diagrams, a clear case where computer typesetting contributes to the aesthetics of the printed page.